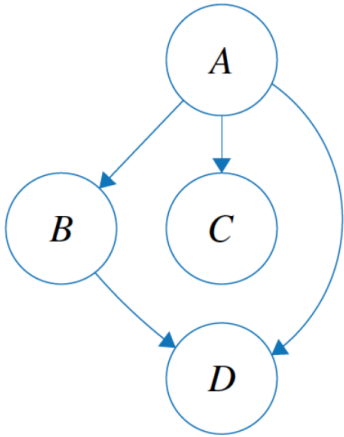


# Graphs Worksheet

## Representation



Draw the adjacency list and the adjacency matrix representation of the above graph.

## DFS Traversal

**Idea:** We must look through all the values of our graph. So, given some starting point, we do a DFS traversal with the caveat that we now track what vertices we've visited before (to avoid cycles)

---

```
public void dfs() {  
    Stack fringe = new Stack();  
    Set visited = new Set();  
    fringe.push(startVertex);  
    while (!fringe.isEmpty()) {  
        Vertex v = fringe.pop();  
        if (!visited.contains(v)) {  
            process(v); //Do something with v  
            for (Vertex neighbor: v.neighbors) {  
                fringe.push(neighbor);  
            }  
            visited.add(v);  
        }  
    }  
}
```

---

# Topological Sort

**Idea:** Given a directed, acyclic graph, how do we 'sort' the vertices based on their dependencies on one another?

---

```
public void topologicalSort() {
    Stack fringe = new Stack();
    Map currentInDegree = new Map<Vertex, Integer>();

    while (!fringe.isEmpty()) {
        Vertex v = fringe.pop();
        process(v); //Do something with v
        for (Vertex neighbor: v.neighbors) {
            currentInDegree(neighbor) -= 1; //Not actual Java code
            if (currentInDegree(neighbor) == 0) {
                fringe.push(neighbor);
            }
        }
    }
}
```

---

**Exercise:** Topologically sort the graph given in class. Draw out the `currentInDegree` map as well as the `fringe`.

## Practice Problems

1. Provide a brief description of how to solve each of the following graph problems efficiently. Provide a worst case runtime bound in  $\Theta$  notation in terms of  $V$  and  $E$ .

(a) Find a path from node  $s$  to node  $t$  in a strongly connected, directed graph.

(b) Determine if a cycle exists in a directed graph.

2. Consider the following implementation of DFS, which contains a crucial error:

---

```
public void dfs() {
    Stack fringe = new Stack();
    Set visited = new Set();

    fringe.add(startVertex);
    visited.add(startVertex);
    while (!fringe.isEmpty()) {
        v = fringe.pop();
        process(v);
        for (Vertex n: v.neighbors()) {
            if (!visited.contains(neighbor)) {
                fringe.push(neighbor);
                visited.add(neighbor);
            }
        }
    }
}
```

---

Give an example of a graph where this algorithm may not traverse in DFS order.

Answers to #1 at <http://datastructur.es/sp17/materials/discussion/discussion11epsol.pdf>

Answers to #2 at <http://datastructur.es/sp17/materials/discussion/discussion11sol.pdf>